COLLECTION OF FORMULAS, TSIN02

1. The distortion due to uniform quantization of a signal sample is

$$D \approx \frac{\Delta^2}{12} = \frac{1}{12} \cdot (\frac{A}{M})^2$$

where Δ is the quantizer step, *A* is the total range of the signal and *M* is the number of levels in the quantizer.

2. General result for the signal-to-distortion ratio for a quantized signal sample.

$$SDR \approx 6R - const$$
 [dB]

where *R* is the number of bits per sample. For uniform quantizer and Gaussian distributed signal, *const*= 7.4.

3. The entropy of a discrete random variable *X* with value set *A*, is

$$H(X) = -\sum_{x \in A} p(x) \cdot \log_2 p(x)$$
 bits

4. For an analog (Gaussian) source with equally strong frequencies within the bandwidth B and a tolerable signal-to-distortion (S/D) ratio, the lowest number of bits per second is

$$R = B\log_2 \frac{S}{D} \approx 0.33B * SDR$$
 (SDR = S/D in dB)

5. Shannon theorem gives an upper bound to the capacity of a link in bits per second (bps), as a function of the available bandwidth and the signal-to-noise ratio of the link

$$C = B \cdot \log_2(1 + \frac{S}{N})$$

where *B* is the bandwidth, *S* is the average signal power and *N* is the average noise power.

6. The Shannon upper bound to the capacity of a binary symmetric channel (BSC)

$$C = 1 - h(p)$$
 [correct bits per transmitted bits]

where *p* is the error probability of the channel and h(p) is the binary entropy function.

$$h(p) = -p \cdot \log_2 p - (1-p) \cdot \log_2 (1-p)$$

7. The packet error probability for a binary symmetric channel (BSC), is

$$P = 1 - (1 - p)^N \approx Np$$
 (if $Np << 1$)

where *N* is the numbers of bits per packet and *p* is the bit error probability.

8. The capacity for the packet loss channel with i.i.d. losses, and packet loss probability P, is

$$C = 1 - P$$

bits/channel symbol.

9. For the first order auto-regressive (AR) model:

$$x(n) = ax_{n-1} + u_n$$

where u_n is zero-mean Gaussian with variance 1, and a < 1,

the average signal power:

$$\sigma^2 = \frac{1}{1 - a^2}$$

the best interpolation between two neighboring samples:

$$\widetilde{x}(n) = \frac{a}{1+a^2} \cdot [x(n-1) + x(n+1)]$$

the average interpolation error:

$$MSE_{interp} = \frac{1}{1+a^2}$$

.

the best prediction of a future sample:

$$\widetilde{x}(n+k) = a^{|k|} \cdot x(n)$$

The average prediction error MSE_{pred}

$$MSE_{pred} = \frac{1-a^{2|k|}}{1-a^2}$$

where k can take on all integers (also negative and 0).

10. The utility function U can be reverse engineered from the demand function D

$$U = \int D^{-1}(p) dp$$

if *U* is a smooth concave function.

10. Optical communication bands: Original band (O-band): 1260 to 1360nm Extended band (E-band): 1360 to 1460nm Short band (S-band): 1460 to 1530nm Conventional band (C-band): 1530 to 1565nm Long band (L-band): 1565 to 1625nm Ultralong band (U-band): 1625 to 1675nm

UDP datagram:

Source port number (16 bits)	Destination port number (16 bits)				
UDP length (16 bits)	UDP checksum (16 bits)				
Data					

TCP segment:

Source port number (16 bits)				Destination port number (16 bits)			
Sequence number (32 bits)							
Acknowledgment (32 bits)							
Header length (4 bits)	Reserved (6 bits)	U A P R C S G K H	R S S Y	5 H Y I N I	N	Window size (16 bits)	
	TCP checksum (16 bits)					Urgent pointer (16 bits)	
Options (if any)							
Data (if any)							

IPv4 datagram:

Version (4 bits)	Header length (4 bits)	Type of Service (8 bits)		Total length (16 bits)		
Identification (16 bits)			Flags (3 bits)	Fragment offset (13 bits)		
Time To L	ive (8 bits)	Protocol (8 bits) †	Checksum (16 bits)			
Source IP address (32 bits)						
Destination IP address (32 bits)						
Options (if any)						
Data						